# The degree/diameter problem in maximal planar bipartite graphs

#### Cristina Dalfó\*, Clemens Huemer\*, Julian Salas\*\*

\*Departament de Matemàtica Aplicada 4 Universitat Politècnica de Catalunya, BarcelonaTech, Barcelona, Catalonia \*\*IIIA-CSIC, Bellaterra, Catalonia

The 6th International Workshop on Optimal Network Topologies Bratislava, Slovakia, June 30 - July 4, 2014

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### Outlook

#### 1. Introduction

2. The  $(\Delta, 2)$  and  $(\Delta, 3)$  problems in maximal planar bipartite graphs

- 3. The  $(\Delta, D)$  problem in maximal planar bipartite graphs
  - 3.1. An upper bound
  - 3.2. A lower bound

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$$|V| - |E| + |F| = 2$$
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 $|E| = 2n - 4$ ,  
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$$|V| - |E| + |F| = 2$$
 (Euler characteristic),  
 $|E| = 2n - 4$ ,  
 $|F| = n - 2$ .

•  $(\Delta, D)$  problem: It consists of finding the maximum possible number of vertices n = |V| in a graph G with maximum degree  $\Delta$  and diameter D.

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# The $(\Delta, 2)$ and $(\Delta, 3)$ problems in maximal planar bipartite graphs

• The  $(\Delta, 2)$  problem:

**Proposition 1.** Consider a maximal planar bipartite graph G with diameter D = 2, maximum degree  $\Delta$  and maximum number of vertices n, then  $n = \Delta + 2$ . The only graph that satisfies this equation is the complete bipartite graph  $K_{2,\Delta}$ .

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# The $(\Delta, 2)$ and $(\Delta, 3)$ problems in maximal planar bipartite graphs

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• The  $(\Delta, 3)$  problem:

**Theorem 2.** Consider a maximal planar bipartite graph G with diameter D = 3, maximum degree  $\Delta$  and maximum number of vertices n, then

$$n = \begin{cases} 3\Delta - 1 & \text{if } \Delta \text{ is odd,} \\ 3\Delta - 2 & \text{if } \Delta \text{ is even.} \end{cases}$$

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• Theorem [Lipton, Tarjan, 1979]. Let G be a planar graph on n vertices containing a spanning tree of radius r. Then V(G) can be partitioned into sets A, B and C such that no edges join vertices in A with vertices in B,  $|A| \leq \frac{2}{3}n$ ,  $|B| \leq \frac{2}{3}n$ , and  $|C| \leq 2r + 1$ .



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 Theorem [Fellows, Hell, Seyffarth, 1995]. Consider a maximal planar graph G with diameter D, maximum degree Δ and maximum number of vertices n, then

$$n = 3(2D+1)(2\Delta^{\lfloor D/2 \rfloor}+1).$$

**Theorem 3.** Let G be a maximal planar bipartite graph on n vertices with maximum degree  $\Delta \ge 4$  and diameter  $D \ge 4$ . Then,

• If  $\Delta = 4$ :  $n \le 6(2D+1)\left(\left\lfloor \frac{D}{2} \right\rfloor^2 + \left\lfloor \frac{D}{2} \right\rfloor + 1\right)$ .

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• If  $\Delta > 4$ :

$$\begin{split} n &\leq 3(2D+1) \left[ \frac{\sqrt{\Delta(\Delta-4)}}{2(\Delta-4)^2} \left[ \left(\Delta - 4 + \sqrt{\Delta(\Delta-4)}\right) \left(\frac{\Delta - 2 - \sqrt{\Delta(\Delta-4)}}{2}\right)^{\lfloor D/2 \rfloor + 1} \right. \\ &\left. - 2\sqrt{\Delta(\Delta-4)} \right. \\ &\left. + \left(4 - \Delta + \sqrt{\Delta(\Delta-4)}\right) \left(\frac{\Delta - 2 + \sqrt{\Delta(\Delta-4)}}{2}\right)^{\lfloor D/2 \rfloor + 1} \right] + 2 \right], \end{split}$$

which is approximately  $3(2D+1)\left[(\Delta-2)^{\lfloor D/2 \rfloor}+1\right]$  if  $\Delta$  is sufficiently large.

#### Proof [sketch].

• We compute from each vertex of C the maximum possible number of vertices at distance at most  $\lfloor D/2 \rfloor$ .

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#### Proof [sketch].

- We compute from each vertex of C the maximum possible number of vertices at distance at most  $\lfloor D/2 \rfloor$ .
- We build a subgraph adding vertices at distance i from a given (root) vertex of C in step i  $(0 \le i \le \lfloor D/2 \rfloor)$ , to obtain an almost maximal (its interior faces are quadrangles) planar bipartite graph.



Figure: An almost maximal subgraph for  $\Delta=4$ 

#### Proof [sketch].

• Let  $n_i$  be the number of vertices at distance i (for  $0 \le i \le \lfloor D/2 \rfloor$ ). For  $i \ge 3$ ,  $n_i$  follows the recurrence

$$n_i = (\Delta - 2)n_{i-1} - n_{i-2}.$$

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#### Proof [sketch].

• Let  $n_i$  be the number of vertices at distance i (for  $0 \le i \le \lfloor D/2 \rfloor$ ). For  $i \ge 3$ ,  $n_i$  follows the recurrence

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• We use the generating function  $G(z) = \frac{\Delta}{\alpha-\beta} \left(\frac{\alpha}{z-\alpha} - \frac{\beta}{z-\beta}\right)$ , where  $\alpha = \frac{1}{2}(\Delta - 2 + \sqrt{\Delta(\Delta - 4)})$  and  $\beta = \frac{1}{2}(\Delta - 2 - \sqrt{\Delta(\Delta - 4)})$  for  $\Delta > 4$ .

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- We obtain for  $\Delta > 4$

$$n_i = \frac{\Delta}{\sqrt{\Delta(\Delta-4)}} \left[ \left( \frac{\Delta - 2 + \sqrt{\Delta(\Delta-4)}}{2} \right)^i - \left( \frac{\Delta - 2 - \sqrt{\Delta(\Delta-4)}}{2} \right)^i \right].$$

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• The total number of vertices  $n = \sum_{i=0}^{\lfloor D/2 \rfloor} n_i$  is obtained as the difference of two geometric series.

• Fellows, Hell, Seyffarth's bound on n for maximal planar graphs:

 $n \leq 3(2D+1)(2\Delta^{\lfloor D/2 \rfloor}+1).$ 

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• Fellows, Hell, Seyffarth's bound on n for maximal planar graphs:

$$n \leq 3(2D+1)(2\Delta^{\lfloor D/2 \rfloor}+1).$$

• D., Huemer, Salas's bound on n for maximal planar bipartite graphs:

$$n \le 3(2D+1)\left[(\Delta-2)^{\lfloor D/2 \rfloor}+1\right].$$

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Figure: Plot of the log (base 10) of the number of vertices n with respect to the diameter D (black points: D., Huemer, Salas's bound; grey points: Fellows, Hell and Seyffarth's bound), for  $\Delta = 5$  and  $4 \le D \le 42$ 

• Ball (of center  $v \in G$  and radius k): it consists of all vertices of G at distance at most k from v.

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- Ball (of center  $v \in G$  and radius k): it consists of all vertices of G at distance at most k from v.
- Theorem [Chepoi, Estellon, and Vaxès]. There exists a constant *C* such that any planar graph *G* of diameter *D* ≤ 2*k* can be covered with at most *C* balls of radius *k*.
- Lower bound: Gavoille, Peleg, Raspaud, and Sopena presented a family of planar graphs with  $C \ge 4$ .

• Corollary 4. There exists a constant C such that each maximal planar bipartite graph G with maximum degree  $\Delta$  and diameter D has at most  $n \leq C(\Delta - 2)^{\lceil D/2 \rceil}$  vertices.

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- For the case D odd and  $\Delta \geq D$ : We use the N-separator theorem by Tishchenko.



Figure: A 5-separator divides the plane into five regions

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 Theorem 5. There exists a constant C such that each maximal planar bipartite graph G with maximum degree Δ and odd diameter D, for Δ ≥ D, has at most n ≤ C(Δ - 2)<sup>LD/2</sup> vertices.

Theorem 6. (a) For any diameter D = 2k (k ≥ 1) and maximum degree Δ (Δ ≥ 5), there exists a maximal planar bipartite graph G<sub>Δ,D</sub> whose number of vertices n(G<sub>Δ,D</sub>) is

$$\frac{\Delta \left(\Delta - 2 + \sqrt{\Delta(\Delta - 4)}\right)^k + \Delta \left(\Delta - 2 - \sqrt{\Delta(\Delta - 4)}\right)^k}{(\Delta - 4)2^k} - \frac{8}{\Delta - 4},$$

which is approximately  $(\Delta - 2)^k$ , for  $\Delta$  and D sufficiently large.



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Figure: The superior half of a maximal planar bipartite graph drawn on a sphere for  $\Delta=4$ 

 Theorem 6. (b) For any diameter D = 2k + 1 (k ≥ 1) and odd maximum degree Δ (Δ ≥ 9), there exists a maximal planar bipartite graph G<sub>Δ,D</sub> whose number of vertices n(G<sub>Δ,D</sub>) is

$$\begin{split} &n(G_{\Delta,3}) = 3\Delta - 1 & \text{for } D = 3, \\ &n(G_{\Delta,5}) = 3\Delta^2 - 21\Delta + 26 & \text{for } D = 5, \\ &n(G_{\Delta,2k+1}) = 3\Delta^2 - 21\Delta + 26 + \frac{3(\Delta - 7)(\Delta - 2)^2((\Delta - 3)^{k-2} - 1)}{(\Delta - 4)} & \text{for } D = 2k + 1 \\ & \text{and } k > 2, \end{split}$$

which is approximately  $3(\Delta - 3)^k$ , for  $\Delta$  and D sufficiently large.

 Theorem 6. (c) For any diameter D = 2k + 1 (k ≥ 1) and even maximum degree Δ (Δ ≥ 10), there exists a maximal planar bipartite graph G<sub>Δ,D</sub> whose number of vertices n(G<sub>Δ,D</sub>) is

$$\begin{split} &n(G_{\Delta,3}) = 3\Delta - 2 & \text{for } D = 3, \\ &n(G_{\Delta,5}) = 3\Delta^2 - 22\Delta + 26 & \text{for } D = 5, \\ &n(G_{\Delta,2k+1}) = 3\Delta^2 - 22\Delta + 26 + \frac{(3\Delta - 22)(\Delta - 2)^2((\Delta - 3)^{k-2} - 1)}{(\Delta - 4)} & \text{for } D = 2k + 1 \\ & \text{and } k > 2, \end{split}$$

which is approximately  $3(\Delta - 3)^k$ , for  $\Delta$  and D sufficiently large.

• Iterative construction for **Theorem 6.** (b), (c):



Figure: The iterative construction

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## Thank you for your attention.

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