

Prisms Aren't Boring

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Joint work with G. Cunningham and M. Mixer



The most important thing...

SIGMAP 2022 is in Fairbanks, Alaska July 10-16.

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UNDER CONSTRUCTION!

Warning

WARNING: Theorems have not been checked for contaminants or adulterants. Theorems have not been peer reviewed and may lead you to draw erroneous conclusions. Theorems are provided without warranty, express or implied.

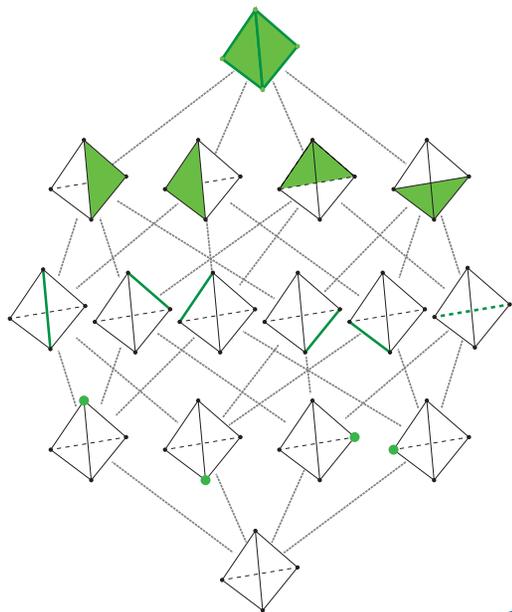


What is a prism?

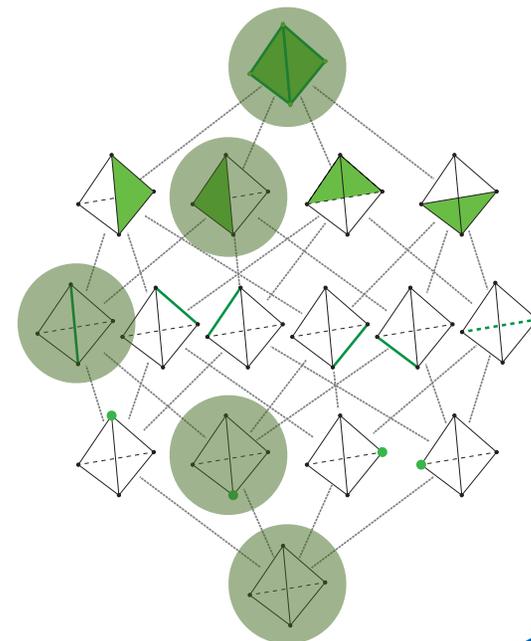
- Geometrically: $\text{Prism}(\Delta) =$ 
- Abstractly:
 - $\text{Prism}(\mathcal{P}) = \mathcal{P} \times \mathcal{J} = \{(f, g) : \mathcal{P}_{-1} \neq f \in \mathcal{P}, \emptyset \neq g \in \mathcal{J}\} \cup \{(\mathcal{P}_{-1}, \emptyset)\}$ where $\mathcal{J} = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$.



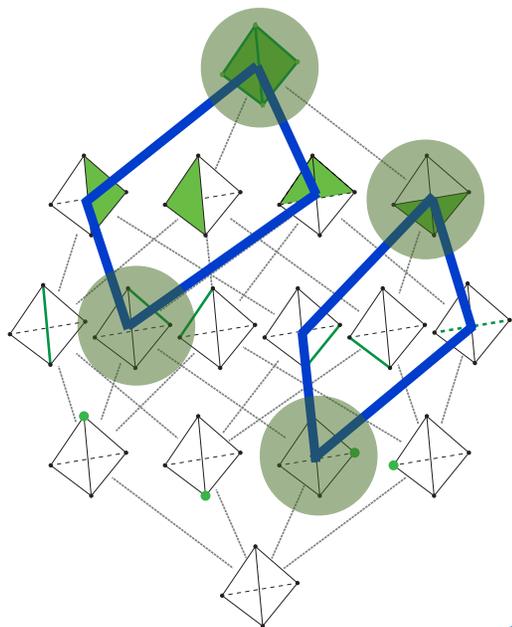
Posets as
Polytopes



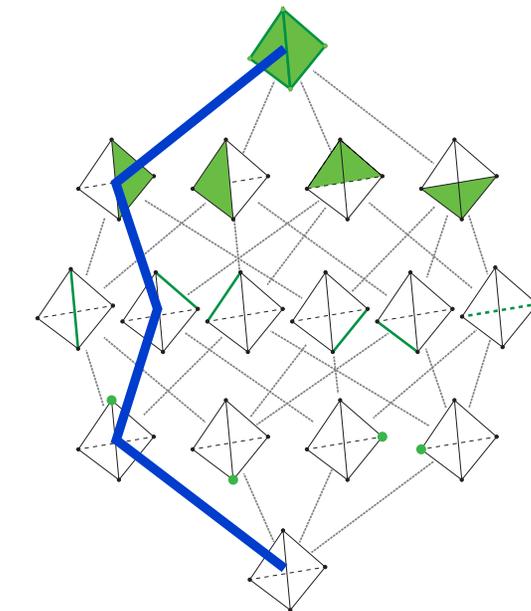
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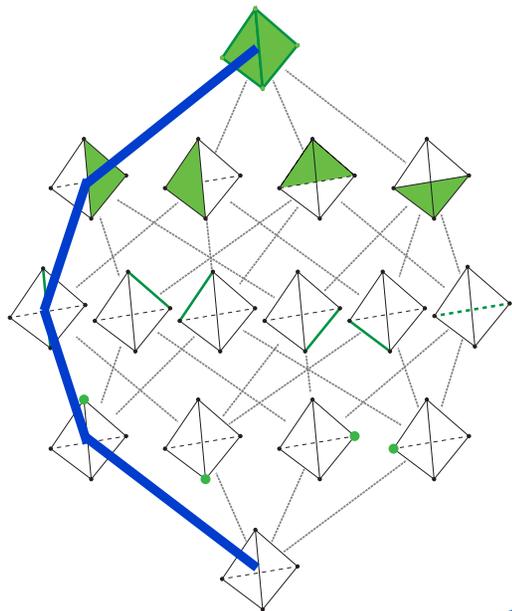
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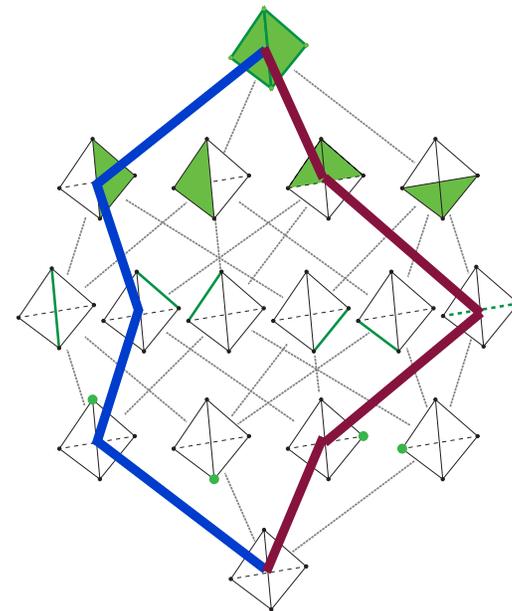
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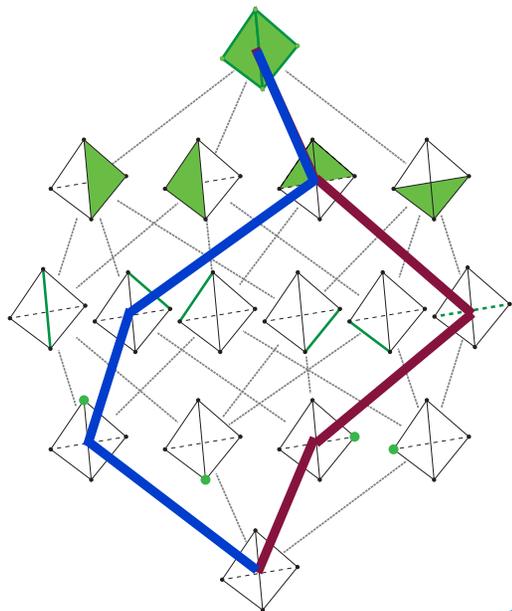
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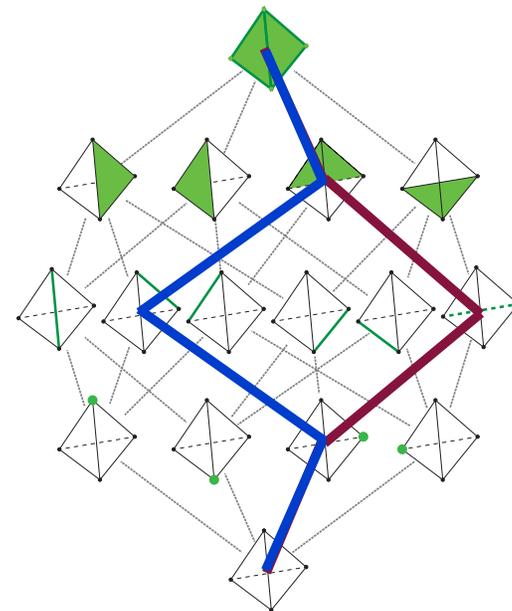
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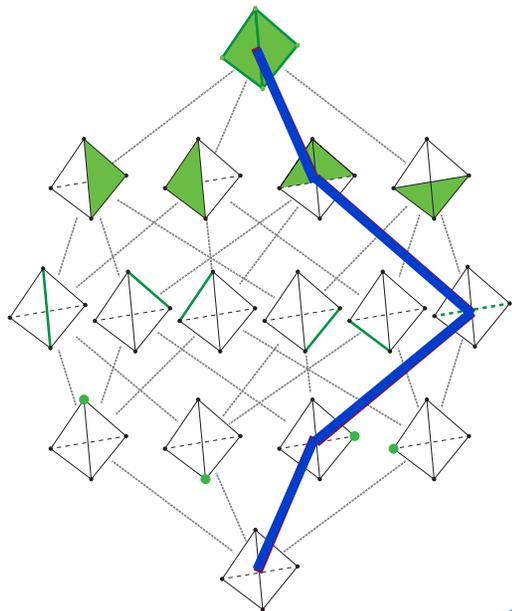
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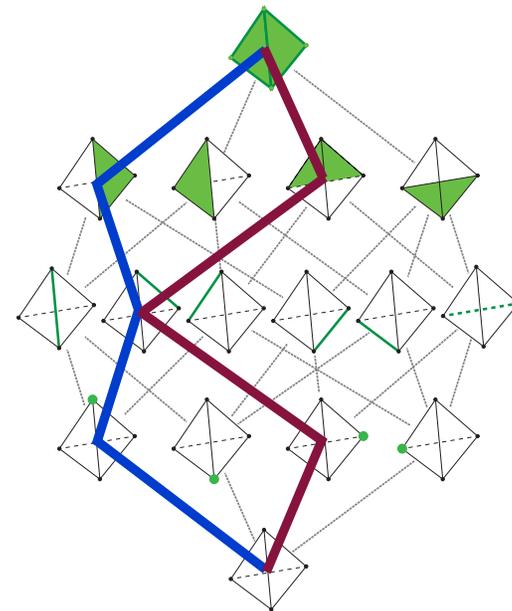
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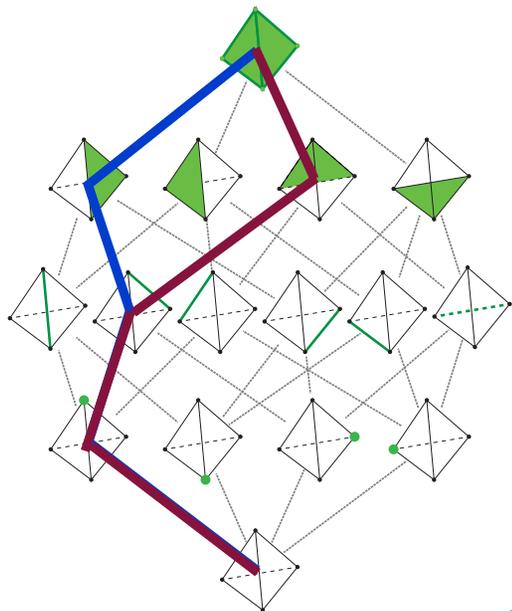
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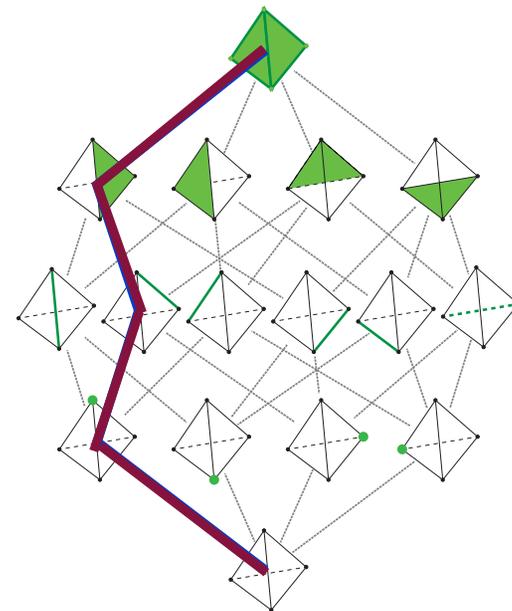
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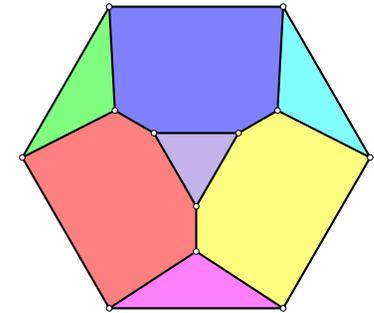
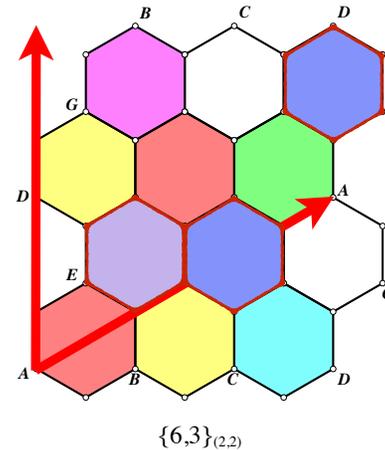


Quick Facts About Abstract Polytopes

- A polytope is *regular* if its automorphism group acts transitively on the flags.
- The *connection group* is $\text{Con}(\mathcal{P}) = \langle r_0, r_1, \dots, r_{d-1} \rangle$
- A polytope \mathcal{P} is a *cover* of a polytope \mathcal{Q} if there exists a rank and adjacency preserving map $\pi : \mathcal{P} \rightarrow \mathcal{Q}$. We write $\mathcal{P} \searrow \mathcal{Q}$.
- A cover \mathcal{R} of a polytope \mathcal{P} is *minimal* if $\mathcal{R} \neq \mathcal{P}$ and $\mathcal{R} \searrow \mathcal{Q} \searrow \mathcal{P}$ implies $\mathcal{R} = \mathcal{Q}$ or $\mathcal{Q} = \mathcal{P}$.



A Regular Cover for the Truncated Tetrahedron



(3.6.6)



Quick Facts About Abstract Polytopes

- A *string group generated by involutions (sggi)* is a group generated by involutions $\rho_0, \rho_1, \dots, \rho_{d-1}$ satisfying $(\rho_j \rho_k)^2 = 1$ if $|j - k| > 1$.
- A *string C-group* is an sggi satisfying the intersection condition:

$$\langle \rho_k \mid k \in I \rangle \cap \langle \rho_k \mid k \in J \rangle = \langle \rho_k \mid k \in I \cap J \rangle$$
- There is a 1-1 correspondence between regular abstract polytopes and string C-groups. The automorphism group of a regular abstract polytope is always a string C-group.



Polytopal Regular Covers

(Monson, Pellicer, W. 2012) The Tomotope has infinitely many distinct minimal regular covers.

(Monson, Pellicer, W. 2014)

- If the connection group of a polytope \mathcal{P} is a string C-group, then \mathcal{P} has a *unique* minimal regular cover.
 - We will say \mathcal{P} is *C-connected*.
- The connection group of every abstract polyhedron is a string C-group, i.e., every polyhedron has a unique minimal regular cover.



Prisms and Covers

Turns out, they aren't boring:

- G. Cunningham, M. Mixer and G.W. have been developing a package library for GAP called **RAMP** (**R**esearch **A**ssistant for **M**aniplexes and **P**olytopes) to make working with maps, maniplexes and abstract polytopes *much* easier.
- We've been using RAMP to develop conjectures.
- There are also LOTS of 4-polytopes \mathcal{P} that **are** C-connected. However...



But...

- Let $\mathcal{P} = \{\{6,3\}_{(2,0)},2\}$, then $\text{Con}(\text{Prism}(\mathcal{P}))$ is **not** a string C-group.
- Let $\mathcal{P} = \{3,3,2\}$, then $\text{Con}(\text{Prism}(\mathcal{P}))$ is **not** a string C-group.
- Let $\mathcal{P} = \{3,3,4\}$, then $\text{Con}(\text{Prism}(\mathcal{P}))$ is **not** a string C-group.
- We also tested more than 1600 small regular polyhedra, and in each case, $\text{Prism}(\mathcal{P})$ was C-connected. Hmm...



In fact...

Main Theorem (Cunningham, Mixer, W.) Let \mathcal{B} be a polyhedron. Then $\text{Prism}(\mathcal{B})$ is C-connected.



Showing a Group is String C

An approach:

(ARP 2E17) Let Γ be an sggi, Δ a string C-group and $\pi : \Gamma \twoheadrightarrow \Delta$, where $\pi(\rho_i) = \sigma_i, \forall i$ and that is one-to-one on $\langle \rho_0, \dots, \rho_{d-2} \rangle$ or $\langle \rho_1, \dots, \rho_{d-1} \rangle$. Then Γ is also a string C-group.

(Pellicer, W. 2018) Let \mathcal{P} be a polytope and $H \leq \text{Aut}(\mathcal{P})$ such that $\mathcal{Q} := \mathcal{P}/H$ is a **pre-maniplex** with the property that $\text{Con}(\mathcal{Q})$ is a string C-group. Let

$L := \{l \in \text{Con}(\mathcal{P}) \mid \forall \Phi \in \mathcal{F}(\mathcal{P}), \exists h \in H \text{ s.t. } l\Phi = \Phi h\}$.

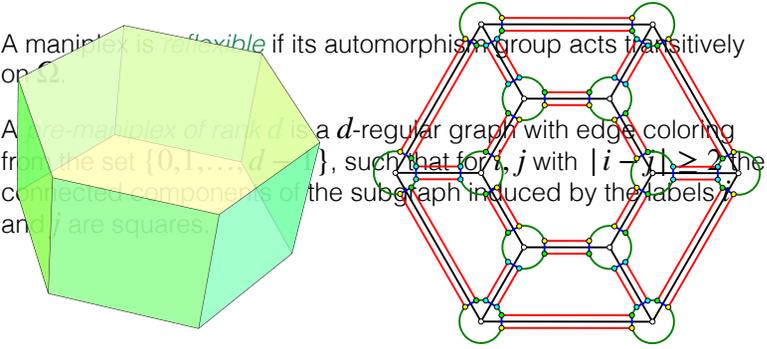
Finally, suppose that if $h \in H$ fixes an incident vertex-facet pair, then $h = 1$. Then $L \triangleleft \text{Con}(\mathcal{P})$, $\text{Con}(\mathcal{P})/L \cong \text{Con}(\mathcal{Q})$, and $\text{Con}(\mathcal{P})$ is a string C-group.



Maniplexes???

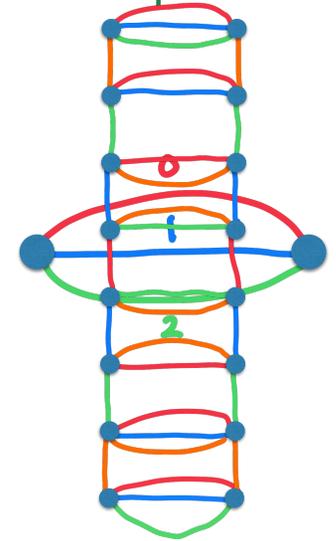
- A *maniplex* \mathcal{M} of rank d is an ordered pair $(\Omega, [r_0, r_1, \dots, r_{d-1}])$ where Ω is a set whose elements are called *flags*, and r_i is a fixed-point-free involution on Ω satisfying
 - $\text{Con}(\mathcal{M}) := \langle r_0, r_1, \dots, r_{d-1} \rangle$ acts transitively on Ω
 - If $|i - j| \geq 2$ then $r_i r_j = r_j r_i$
 - If $i \neq j$ then r_i and r_j have no transpositions in common

- A maniplex is *reflexible* if its automorphism group acts transitively on Ω .
- A *pre-maniplex of rank d* is a d -regular graph with edge coloring from the set $\{0, 1, \dots, d-1\}$, such that for i, j with $|i - j| \geq 2$ the connected components of the subgraph induced by the labels i and j are squares.



A Useful Pre-Maniplex

- \mathcal{G}_3 is the pre-maniplex with two flags of rank 3.
- $\mathcal{Q} := \text{Prism}(\mathcal{G}_3)$
- $\text{Con}(\mathcal{Q})$ is a string C-group



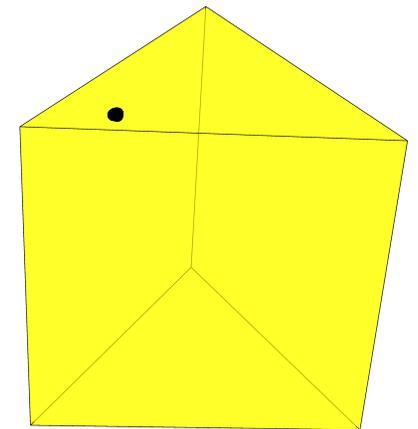
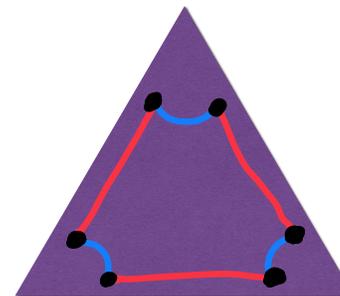
Stratified Operations

(From Cunningham, Pellicer, W.)

Let M_d be the set of all maniplexes of rank d .

An operation $F : M_n \rightarrow M_m$ is *stratified* if there is a set A (the *strata*) such that

- If $\text{Flags}(\mathcal{M}) = \Omega$, then $\text{Flags}(F(\mathcal{M})) \subseteq A \times \Omega$ s.t. the canonical projections are surjective.
- The universal string Coxeter group W_m has a well defined action on A .
- The action of W_m can be described *nicely* in terms of the natural action of W_n on Ω and the action of W_m on A .



Stratified Operations

A stratified operation F is *fully* stratified if $\text{Flags}(F(\mathcal{M})) = A \times \Omega$.

It is *cover-preserving* if, whenever $\mathcal{M} \searrow \mathcal{L}$, then $F(\mathcal{M}) \searrow F(\mathcal{L})$.

Proposition (Cunningham, Pellicer, W.) Fully-stratified operations are cover-preserving.

Proposition (CPW) Let \mathcal{R} be a minimal regular cover of \mathcal{P} , then $\text{Con}(F(\mathcal{R})) \cong \text{Con}(F(\mathcal{P}))$.



Proof for Prisms Over Polyhedra

We show the following:

1. Let \mathcal{B} be an orientable regular polyhedron, then $\text{Con}(\text{Prism}(\mathcal{B}))$ is a string C-group.
 - Proof uses fact $\text{Prism}(\mathcal{B}) \searrow \text{Prism}(\mathcal{G}_3)$.
2. Let \mathcal{M} be a non-orientable reflexible 3-manifold and let \mathcal{L} be its orientable double cover. Then the prism with base \mathcal{L} is the same as the orientable double cover of the prism with base \mathcal{M} .
 - Proof uses fact that ODC is cover preserving.



Proof Steps

3. Recall (CunPelWil) that if F is a parallel product then F is cover-preserving and connection-preserving.

- Also: $\text{Con}(\mathcal{M}) \cong \text{Con}(\text{src}(\mathcal{M}))$

4. Show that the orientable cover operation is a parallel product operation on manifolds.



Proof Steps

5. Show that if \mathcal{M} is a non-orientable 3-manifold and $\mathcal{L} = \text{odc}(\mathcal{M})$, then $\text{src}(\text{Prism}(\mathcal{M})) = \text{src}(\text{Prism}(\mathcal{L}))$.
6. Complete the proof:
 - We have the case when \mathcal{B} is orientable and reflexible.
 - By 5, holds for all reflexible polyhedra.
 - Note $\text{Con}(\mathcal{B})$ is a string C-group for any polyhedron.
 - Use result about Prism being a parallel product to generalize to all polyhedra.



Something Interesting About Cubes

- **Observation 1:** Let \mathcal{B} be a regular polyhedron with automorphism group $\Gamma = \langle \rho_0, \rho_1, \rho_2 \rangle$, and let $\mathcal{P} = \text{Prism}(\mathcal{B})$. Then we may represent $\text{Con}(\mathcal{P}) = \langle s_0, s_1, s_2, s_3 \rangle \leq \mathcal{S}_8 \wr \text{Aut}(\mathcal{B})$ with generators:
 - $s_0 := ((4,5), [\rho_0, \rho_0, \rho_0, e, e, \rho_0, \rho_0, \rho_0])$
 - $s_1 := ((3,4)(5,6), [\rho_1, \rho_1, e, e, e, e, \rho_1, \rho_1])$
 - $s_2 := ((2,3)(6,7), [\rho_2, e, e, \rho_1, \rho_1, e, e, \rho_2])$
 - $s_3 := ((1,2)(7,8), [e, e, \rho_2, \rho_2, \rho_2, \rho_2, e, e])$



Something Interesting About Cubes

- The interesting bit:
 - $(s_0 s_1)^4 = ((), [(\rho_0 \rho_1)^4, (\rho_0 \rho_1)^4, e, e, e, e, (\rho_0 \rho_1)^4, (\rho_0 \rho_1)^4])$
 - $(s_1 s_2)^3 = ((), [(\rho_1 \rho_2)^3, e, e, e, e, e, e, (\rho_1 \rho_2)^3])$
 - $(s_2 s_3)^3 = ((), [e, e, e, (\rho_1 \rho_2)^3, (\rho_1 \rho_2)^3, e, e, e])$
- So if $|(\rho_0 \rho_1)| = 4$ and $|(\rho_1 \rho_2)| = 3$ then \mathcal{P} is the 4-cube (γ_4 in Coxeter's notation).



Something About Cubes

- **Observation 2:** Let \mathcal{G}_d be the pre-maniplex with two flags of rank d . Then we can show that $\text{Con}(\text{Prism}(\mathcal{G}_d)) = \text{Con}(\gamma_d)$.
- **Observation 3:** Let \mathcal{B} be a polyhedron, then $\text{Con}(\mathcal{B} \diamond \gamma_3)$ **is** a string C-group.
- **Observation 4:** (Theorem) Let \mathcal{M} be a reflexible 3-maniplex, and let $\mathcal{P} = \text{Prism}(\mathcal{M})$. Then $\text{src}(\mathcal{P}) \searrow \gamma_4$ and $\text{src}(\mathcal{P}) \diamond \gamma_4 = \text{src}(\mathcal{P})$.



Something About Cubes

- **Observation 5:** For every *examined* \mathcal{B} of rank > 3 : if $\text{Con}(\text{Prism}(\mathcal{B}))$ is **not** a string C-group *then* $\text{Con}(\mathcal{B} \diamond \gamma_d)$ is not a string C-group!

Conjecture: Let \mathcal{B} be an abstract polytope of rank d , and let γ_d be the d -cube. Then $\text{Con}(\text{Prism}(\mathcal{B}))$ is a string C-group iff $\text{Con}(\mathcal{B} \diamond \gamma_d)$ is a string C-group.



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Don't forget the most important thing...

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SIGMAP 2022

Symmetries in Graphs, Maps, and Polytopes Workshop 2022



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General Information

Monday July 10th to Friday July 16th, 2022.

The 2022 Edition of the Symmetries in Graphs, Maps, and Polytopes Workshop will be held at the University of Alaska Fairbanks. There will be nine plenary talks, a problem session, and a number of shorter contributed talks. Time will be set aside in the schedule to encourage and accommodate collaboration for both ongoing research or for starting new projects.

You may pre-register using [this form](#), and we will notify you as soon as the full registration system is available. Information you provide in your pre-registration form will also aid us in advance planning and help us keep costs down.

